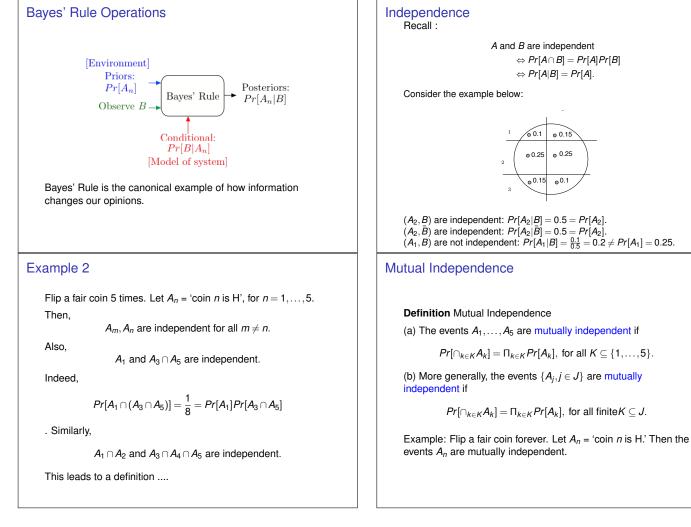
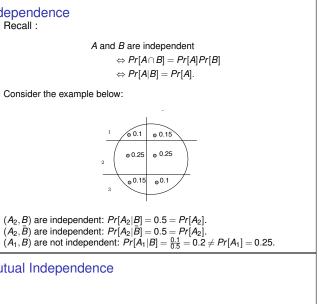


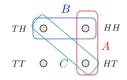
pprox 0.42





Pairwise Independence

- Flip two fair coins. Let
- A = 'first coin is H' = {HT, HH};
- B = 'second coin is H' = {*TH*, *HH*};
- C = 'the two coins are different' = {TH, HT }.



A, C are independent; B, C are independent;  $A \cap B$ , *C* are not independent.  $(Pr[A \cap B \cap C] = 0 \neq Pr[A \cap B]Pr[C])$ .

A did not say anything about C and B did not say anything about C, but  $A \cap B$  said something about C!

Mutual Independence

#### Theorem

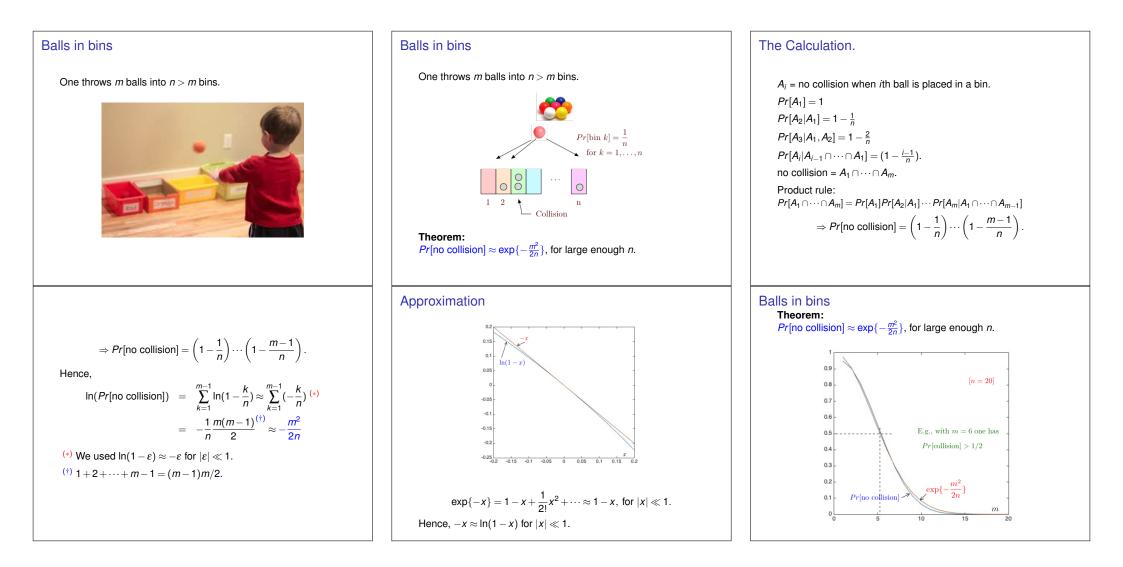
(a) If the events  $\{A_i, j \in J\}$  are mutually independent and if  $K_1$ and  $K_2$  are disjoint finite subsets of J, then

 $\cap_{k \in K_1} A_k$  and  $\cap_{k \in K_2} A_k$  are independent.

(b) More generally, if the  $K_n$  are pairwise disjoint finite subsets of *J*, then the events

 $\cap_{k \in K_n} A_k$  are mutually independent.

(c) Also, the same is true if we replace some of the  $A_k$  by  $\bar{A}_k$ .



Balls in bins	The birthday paradox	Today's your birthday, it's my birthday too
Theorem: $Pr[no collision] \approx exp\{-\frac{m^2}{2n}\}$ , for large enough $n$ .In particular, $Pr[no collision] \approx 1/2$ for $m^2/(2n) \approx \ln(2)$ , i.e., $m \approx \sqrt{2\ln(2)n} \approx 1.2\sqrt{n}$ .E.g., $1.2\sqrt{20} \approx 5.4$ .Roughly, $Pr[collision] \approx 1/2$ for $m = \sqrt{n}$ . $(e^{-0.5} \approx 0.6.)$		Probability that <i>m</i> people all have different birthdays? With <i>n</i> = 365, one finds <i>Pr</i> [collision] $\approx 1/2$ if $m \approx 1.2\sqrt{365} \approx 23$ . If <i>m</i> = 60, we find that <i>Pr</i> [no collision] $\approx \exp\{-\frac{m^2}{2n}\} = \exp\{-\frac{60^2}{2 \times 365}\} \approx 0.007$ . If <i>m</i> = 366, then <i>Pr</i> [no collision] = 0. (No approximation here!)
n p(n)   1 0.0%   5 2.7%   10 11.7%   20 41.1%   23 50.7%   30 70.6%   40 89.1%   50 97.0%   60 99.4%   70 99.99997%   200 99.9999999999999999999999999999999999	$\label{eq:consider} \begin{array}{l} \textbf{Checksums!} \\ \textbf{Consider a set of $m$ files.} \\ \textbf{Each file has a checksum of $b$ bits.} \\ \textbf{How large should $b$ be for $Pr$[share a checksum] $\leq 10^{-3}$?} \\ \textbf{Claim: $b \geq 2.9 \ln(m) + 9$.} \\ \textbf{Proof:} \\ \textbf{Let $n = 2^b$ be the number of checksums.} \\ \textbf{We know $Pr$[no collision] $\approx \exp\{-m^2/(2n)\} $\approx 1 - m^2/(2n)$.} \\ \textbf{Hence,} \\ \begin{array}{l} Pr[\text{no collision}] $\approx 1 - 10^{-3} \Leftrightarrow m^2/(2n) \approx 10^{-3} \\ $\approx 2n \approx m^2 10^3 \Leftrightarrow 2^{b+1} \approx m^2 2^{10} \\ $\Leftrightarrow b + 1 \approx 10 + 2 \log_2(m) \approx 10 + 2.9 \ln(m)$.} \\ \textbf{Note: } \log_2(x) = \log_2(e) \ln(x) \approx 1.44 \ln(x). \end{array}$	Coupon Collector Problem.There are <i>n</i> different baseball cards. (Brian Wilson, Jackie Robinson, Roger Hornsby,) One random baseball card in each cereal box.Theorem: If you buy <i>m</i> boxes, (a) $Pr[miss one specific item] \approx e^{-\frac{m}{n}}$ (b) $Pr[miss any one of the items] \leq ne^{-\frac{m}{n}}$

# Coupon Collector Problem: Analysis.

Event  $A_m$  = 'fail to get Brian Wilson in *m* cereal boxes' Fail the first time:  $(1 - \frac{1}{n})$ Fail the second time:  $(1 - \frac{1}{n})$ And so on ... for *m* times. Hence,

$$Pr[A_m] = (1 - \frac{1}{n}) \times \dots \times (1 - \frac{1}{n})$$
$$= (1 - \frac{1}{n})^m$$
$$ln(Pr[A_m]) = mln(1 - \frac{1}{n}) \approx m \times (-\frac{1}{n})$$
$$Pr[A_m] \approx exp\{-\frac{m}{n}\}.$$

For  $p_m = \frac{1}{2}$ , we need around  $n \ln 2 \approx 0.69n$  boxes.

### Summary.

Bayes' Rule, Mutual Independence, Collisions and Collecting

#### Main results:

- Bayes' Rule:  $Pr[A_m|B] = p_m q_m / (p_1 q_1 + \dots + p_M q_M).$
- ▶ Product Rule:  $Pr[A_1 \cap \cdots \cap A_n] = Pr[A_1]Pr[A_2|A_1] \cdots Pr[A_n|A_1 \cap \cdots \cap A_{n-1}].$
- **Balls in bins**: *m* balls into n > m bins.

$$Pr[\text{no collisions}] \approx \exp\{-\frac{m^2}{2n}\}$$

► Coupon Collection: *n* items. Buy *m* cereal boxes.

 $Pr[\text{miss one specific item}] \approx e^{-\frac{m}{n}}; Pr[\text{miss any one of the items}] \leq ne^{-\frac{m}{n}}.$ 

Key Mathematical Fact:  $\ln(1-\varepsilon) \approx -\varepsilon$ .

## Collect all cards?

Experiment: Choose *m* cards at random with replacement. Events:  $E_k$  = 'fail to get player k', for k = 1, ..., n Probability of failing to get at least one of these *n* players:

 $p := Pr[E_1 \cup E_2 \cdots \cup E_n]$ How does one estimate *p*? Union Bound:  $p = Pr[E_1 \cup E_2 \cdots \cup E_n] \le Pr[E_1] + Pr[E_2] \cdots Pr[E_n].$ 

$$Pr[E_k] \approx e^{-\frac{m}{n}}, k = 1, \ldots, n.$$

Plug in and get

 $p < ne^{-\frac{m}{n}}$ .

## Collect all cards?

Thus,

 $Pr[missing at least one card] \leq ne^{-\frac{m}{n}}.$ 

Hence,

Pr[missing at least one card $] \le p$  when  $m \ge n \ln(\frac{n}{p})$ .

To get p = 1/2, set  $m = n \ln (2n)$ . E.g.,  $n = 10^2 \Rightarrow m = 530$ ;  $n = 10^3 \Rightarrow m = 7600$ .